

Quantum Signal Processing and Quantum Singular Value Transformation on U(N)



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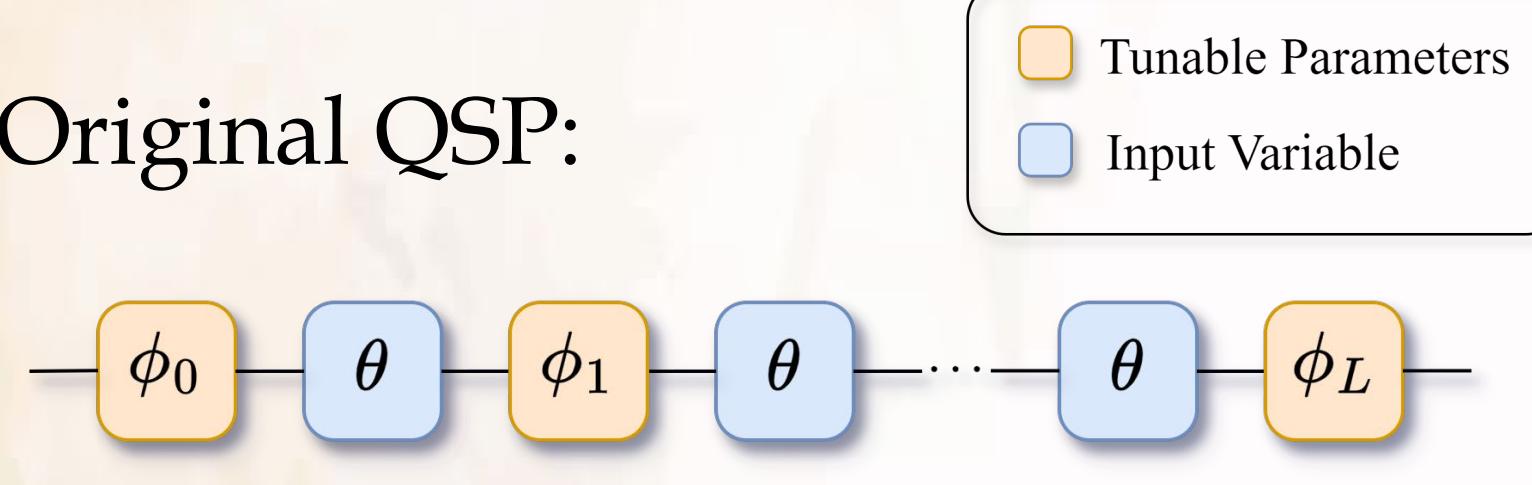


arXiv:2408.01439

Background: U(2)-QSP

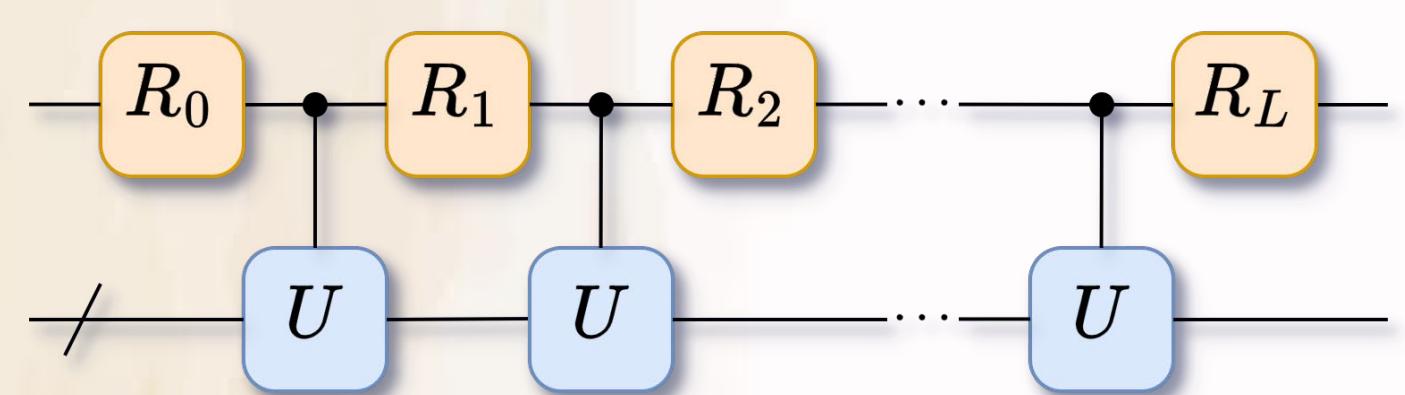
(Quantum Signal Processing)

Original QSP:



- θ : Input signal
- The blue and orange gates are rotation gates around two orthogonal axis (e.g. R_x and R_z)

Generalized QSP:



- U : Input signal
- R_0, \dots, R_L : Arbitrary single-qubit operations

In generalized QSP, there exists $R_1, \dots, R_L \in U(2)$ such that the circuit implements the transfrm:

$$\begin{pmatrix} f(U) & * \\ * & * \end{pmatrix}$$

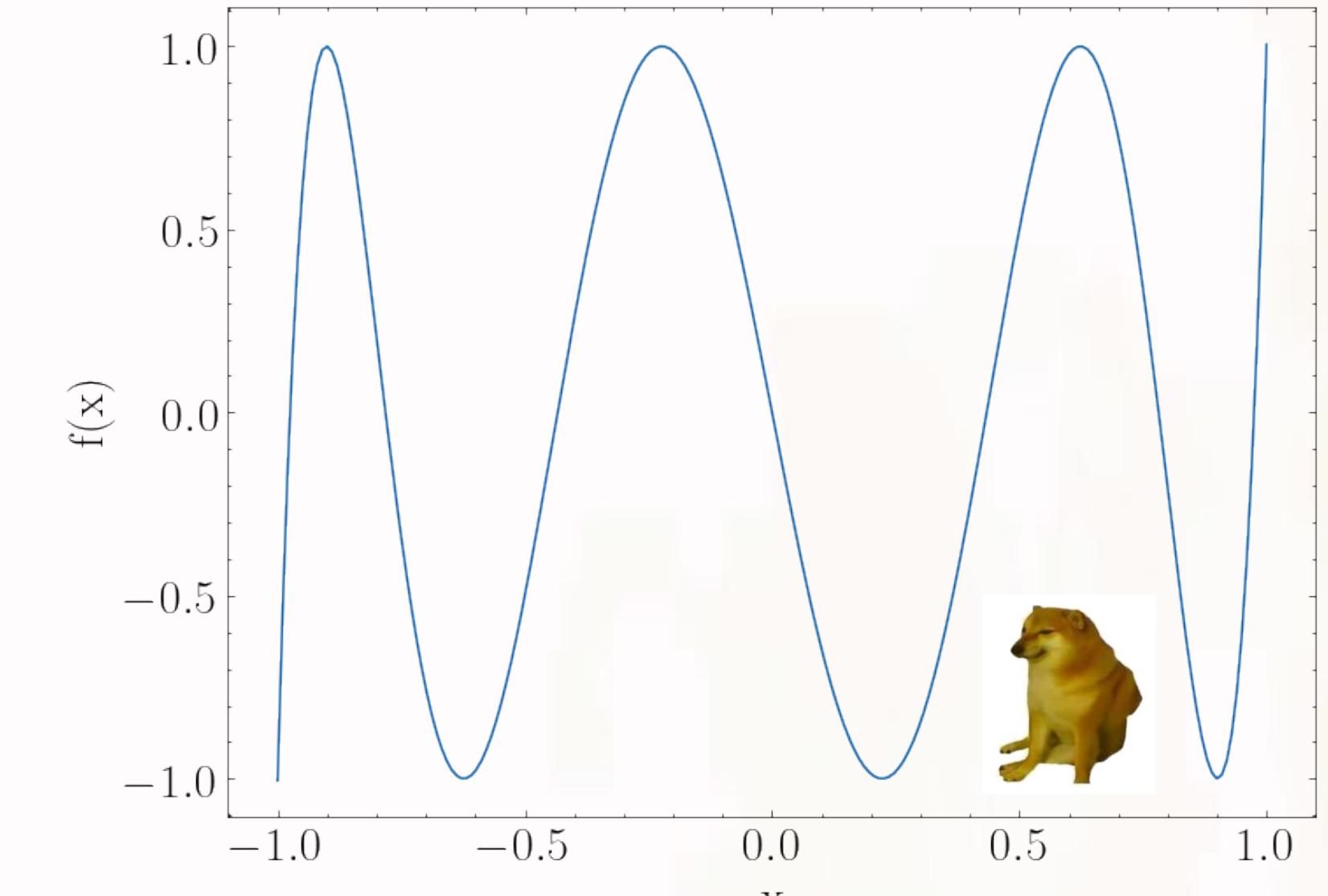
if and only if:

- f is a polynomial, $\deg(f) \leq \# \text{Calls to } U$;
- $|f(z)| \leq 1$ for all $|z| = 1$.

Example applications:

- Fixed-point Search: $f(x) \sim \text{sgn}(x)$;
- Linear System Solver: $f(x) \sim 1/x$;
- Hamiltonian Simulation: $f(x) \sim e^{-itx}$.

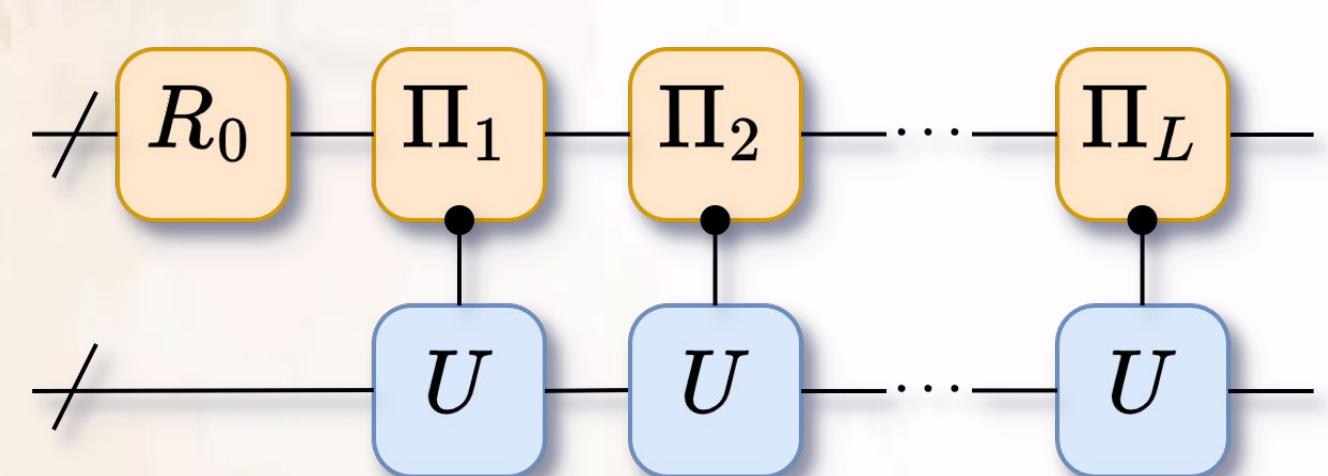
Complexity of using polynomial to approach analytic function: $d = O(\log(\frac{1}{\epsilon}))$



In U(2)-QSP/QSVT, one can only tune the circuit to realize 1 desired polynomial.
E.g.: amplitude amplification realizes the Chebyshev polynomial $\cos(\theta) \mapsto \cos((2k+1)\theta)$.

Main Result: U(N)-QSP

U(N)-QSP:



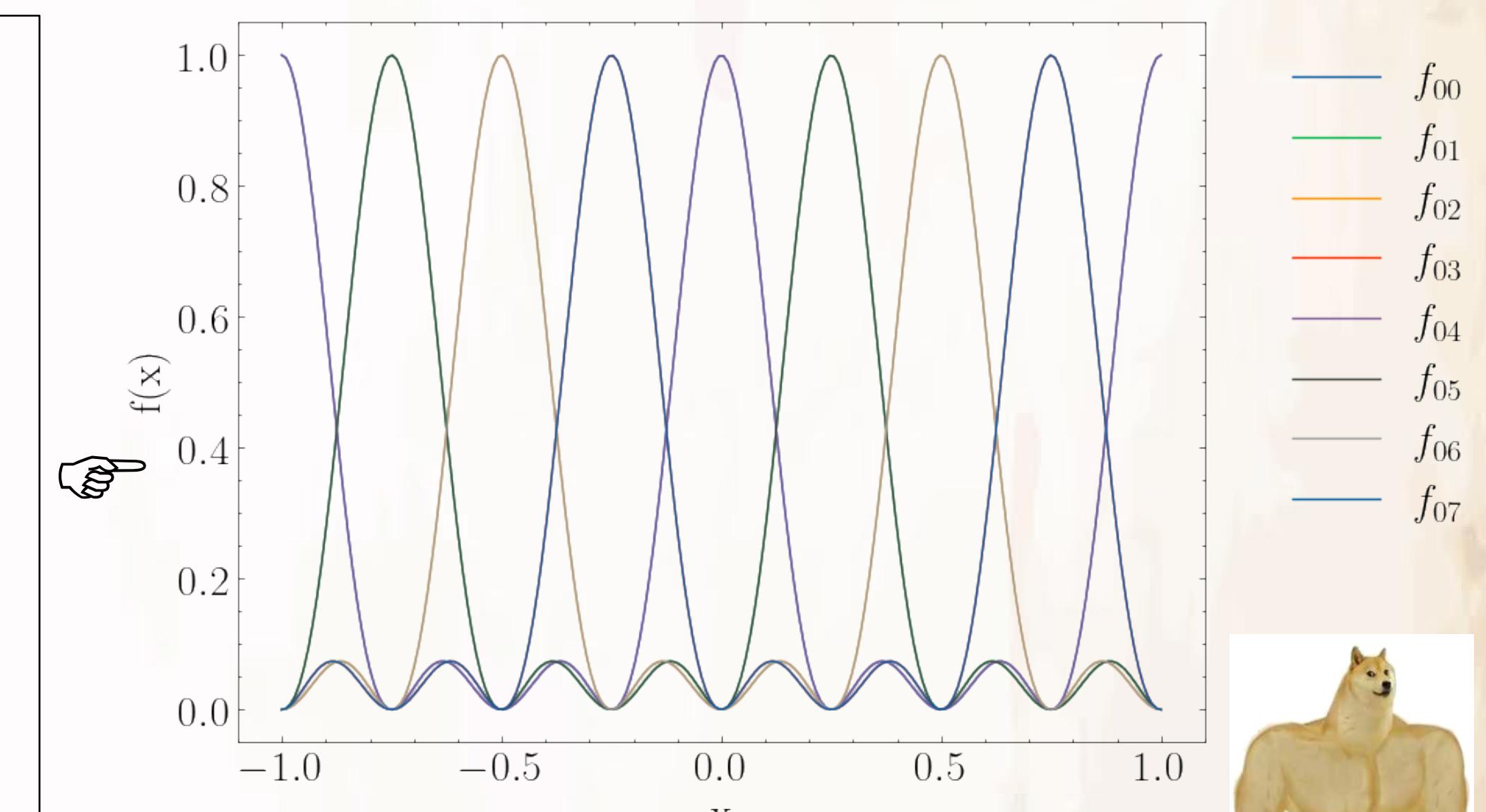
- U : Input signal
- R_0 : tunable n -qubit unitary
- Π_1, \dots : Tunable projector: $C_\Pi(U) = \Pi \otimes U + (I - \Pi) \otimes I$

In $U(N)$ -QSP, there exists $R_0 \in U(N)$ and projectors Π_1, \dots, Π_L such that the circuit implements the transfrm:

$$\begin{pmatrix} f_{00}(U) & f_{01}(U) & \dots & * \\ f_{10}(U) & f_{11}(U) & \dots & * \\ \vdots & \vdots & \ddots & \vdots \\ * & * & \dots & * \end{pmatrix}$$

if and only if:

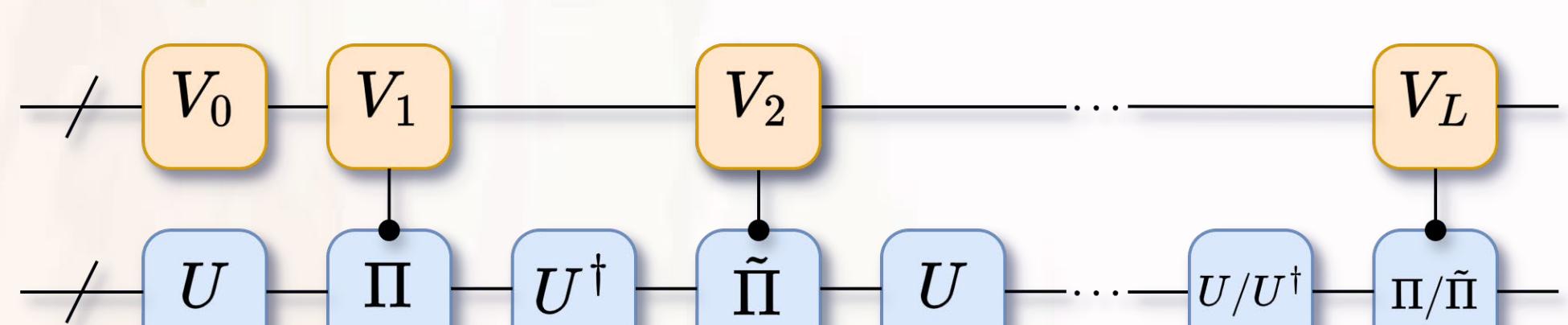
- Each f_{jk} is a polynomial, $\deg(f_{jk}) \leq \# \text{Calls to } U$;
- The matrix $\{f_{jk}(z)\}$ has singular value ≤ 1 for all $|z| = 1$.



In $U(N)$ -QSP/QSVT one can have control of multiple polynomials simultaneously.

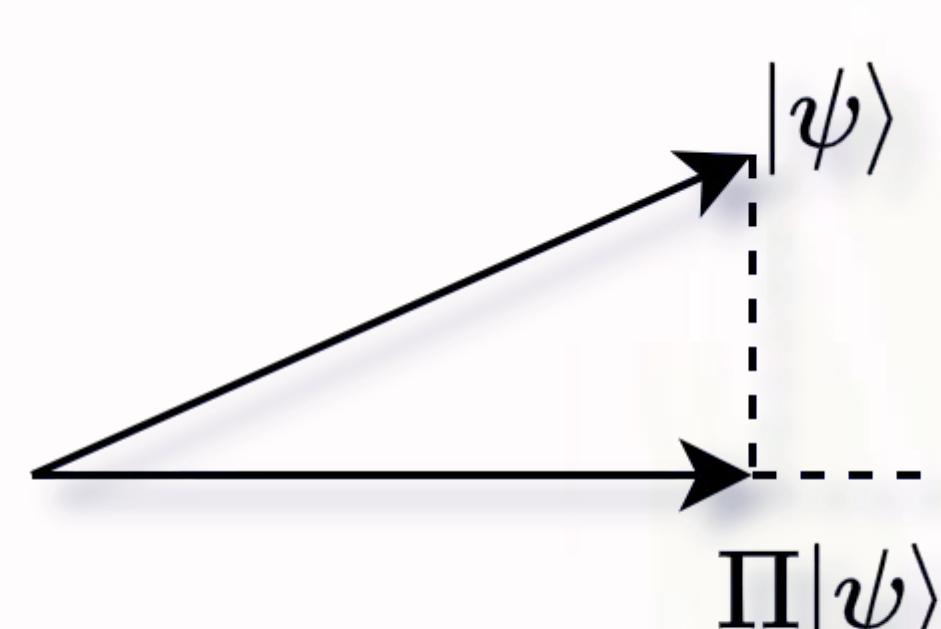
U(N)-QSVT

(Quantum Singular Value Transformation)



- $U = \begin{pmatrix} A & * \\ * & * \end{pmatrix}$: Block encoding of the input signal A
- V_0, \dots, V_L : Tunable $U(N)$ operators
- Output: block encoding of multiple singular-value polynomial transformations

$$f^{(sv)} \left(\sum_j \lambda_j |\psi_j\rangle\langle\phi_j| \right) = \begin{cases} \sum_j f(\lambda_j) |\psi_j\rangle\langle\phi_j|, & f \text{ is even} \\ \sum_j f(\lambda_j) |\phi_j\rangle\langle\psi_j|, & f \text{ is odd} \end{cases}$$



Optimal Quantum Amplitude Estimation:
Estimate $x = \|\Pi|\psi\rangle\|^2$ with optimized number of state preparation oracle of $|\psi\rangle$.

$$\begin{array}{c} \begin{bmatrix} p_0(\hat{w}) & p_1(\hat{w}) & \dots & * \\ \vdots & \vdots & \ddots & \vdots \\ * & * & \dots & * \end{bmatrix} \\ \times \begin{bmatrix} q_0(\hat{v}) & \dots & * \\ q_1(\hat{v}) & \dots & * \\ \vdots & \ddots & \vdots \\ * & \dots & * \end{bmatrix} \end{array} \xrightarrow{\text{Product of block-encoded matrices}} \begin{bmatrix} \sum_j p_j(\hat{w}) q_j(\hat{v}) & * & * \\ * & * & * \end{bmatrix}$$

Combining with other means of manipulating block-encoded matrices, like linear combination and product, $U(N)$ -QSP could provide new approach to general multi-variate QSP[Quantum 6, 811 (2022)].

Open Questions

Can we...

- Find efficient classical algorithm to evaluate/compile R_0, Π_1, \dots ?
- Solve more sensing / metrology problems with $U(N)$ -QSP/QSVT?
- Use $U(N)$ -QSP/QSVT as ansatz in variational quantum algorithms / state preparation?
- Design algorithms on hybrid qubit-oscillator systems with $U(N)$ -QSP/QSVT?
- Construct circuits for general multi-variate QSP with $U(N)$ -QSP/QSVT being a building block?